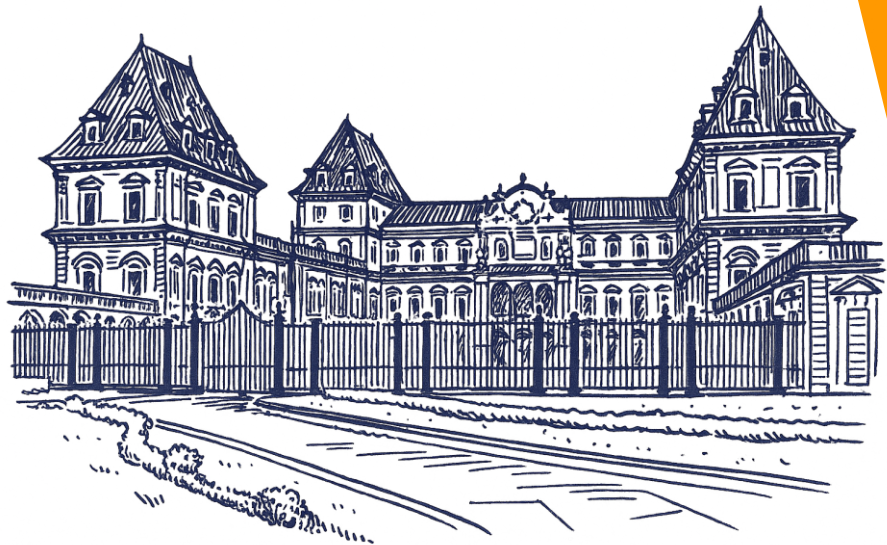


# An aperitif on modern cryptography

Leonardo Errati  
*April 10<sup>th</sup>, 2026*



Politecnico  
di Torino

**CrypTO**

**INTERLUDE: What Cryptography is Not**

# Cryptography?

S S M **W O R D** K N U S  
E W **S E A R C H** J **J** S  
T P S H B D S A S O J  
**E** **F** S G **L O V E** F U K  
N N U C R S E E S R T  
J P S N K E P S E N T  
O D L S C A L L V E C  
Y R S A C K B A X Y H  
D T P S Y X L O X E P  
S E **E P U Z Z L E** S F

NOT cryptography

AF104CE57016B84840AC04  
1E383AC472105A50E0095E  
50467ACD0BD49415076A02  
9828A623D4932111424084  
8BB5190A46630D84630557  
42898C771AC907892E40E7  
965A4DFF46004C89BD7217  
6859DB223061A141C274FE  
AC10700563391DEB92CE06  
019507FB59C34AE145589A

cryptography

# Cryptography?

<b>Cryptography</b>	Making ciphers
<b>Cryptanalysis</b>	Breaking ciphers
<b>Cryptology</b>	Making & breaking ciphers

**1.** Against entropy...

# Atbash cipher

*“And after all of them, the king of  
**Sheshak** will drink it too.”*

---

(Book of Jeremiah, 25:26)

*“Behold, I will raise up against  
Babylon, and against the inhabitants  
of **Lev-kamai**, a destroying wind.”*

---

(Book of Jeremiah, 51:1)

# Atbash cipher

*“And after all of them, the king of  
**Sheshak** will drink it too.”*

---

(Book of Jeremiah, 25:26)

***Sheshak***

*shin-shin-kaf*

# Atbash cipher

“And after all of them, the king of **Sheshak** will drink it too.”

(Book of Jeremiah, 25:26)

**Sheshak**  
shin-shin-kaf  $\xrightarrow{\varphi_{\text{atbash}}}$  bet-bet-lamed

22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1
ת	ש	ר	ק	צ	פ	ע	ס	נ	מ	ל	כ	י	ט	ח	ז	ו	ה	ד	ג	ב	א
Tav	Shin	Reish	Qof	Tsadé	Pé	Ayin	Samekh	Noun	Meme	Lamed	Kaf	Yod	Teith	Heith	Zayin	Vav	Hé	Dalet	Guimel	Beith	Alph

$$\varphi_{\text{atbash}}: \mathbb{Z}_{22} \rightarrow \mathbb{Z}_{22}$$
$$x \mapsto (23 - x) \text{ mod } 22$$

A  $\longrightarrow$  Z  
B  $\longrightarrow$  Y  
C  $\longrightarrow$  X

# Atbash cipher

“And after all of them, the king of **Sheshak** will drink it too.”

(Book of Jeremiah, 25:26)

**Sheshak**  
shin-shin-kaf

$\varphi_{\text{atbash}}$

**Babylon**  
bet-bet-lamed

22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1
ת	ש	ר	ק	צ	פ	ע	ס	נ	מ	ל	כ	י	ט	ח	ז	ו	ה	ד	ג	ב	א
Tav	Shin	Reish	Qof	Tsadé	Pé	Ayin	Samekh	Noun	Meme	Lamed	Kaf	Yod	Teith	Heith	Zayin	Vav	Hé	Dalet	Guimel	Beith	Alph

$$\varphi_{\text{atbash}}: \mathbb{Z}_{22} \rightarrow \mathbb{Z}_{22}$$

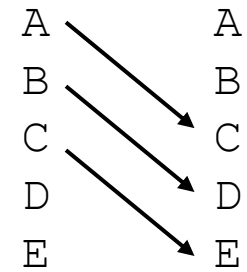
$$x \mapsto (23 - x) \text{ mod } 22$$

A → Z  
B → Y  
C → X

# Caesar cipher



message:       ATTACKATDAWN  
key:            C  
ciphertext:    DWWDFNDWGDZQ



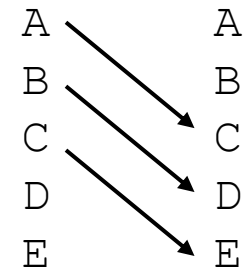
# Caesar cipher



$$\varphi_{\text{caesar}}: \mathbb{Z}_{26} \rightarrow \mathbb{Z}_{26}$$

$$x \mapsto (x + k) \text{ mod } 26$$

message:           ATTACKATDAWN  
key:                C  
ciphertext:        DWWDFNDWGDZQ



# Caesar cipher

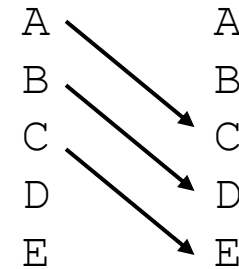
$$H(K) = - \sum_{k \in K} P[k] \cdot \log_2 P[k]$$

$H_{caesar} \approx 4.7$



$$\varphi_{caesar}: \mathbb{Z}_{26} \rightarrow \mathbb{Z}_{26}$$
$$x \mapsto (x + k) \text{ mod } 26$$

message:           ATTACKATDAWN  
key:                C  
ciphertext:        DWWDFNDWGDZQ



# Vigenere cipher

$$k = (k_0, \dots, k_{t-1})$$

$$H_{vigenere} \simeq 4.7 \cdot t$$

$$\varphi_{vigenere}: \mathbb{Z}_{26} \rightarrow \mathbb{Z}_{26}$$

$$x \mapsto (x_i + k_{i \bmod t}) \bmod 26$$

message:	ATTACKATDAWN
key:	keykeykeykey
ciphertext:	KXRKGIKXBKAL



# Vigenere cipher

$$k = (k_0, \dots, k_{t-1})$$

$$H_{\text{vigenere}} \simeq 4.7 \cdot t$$

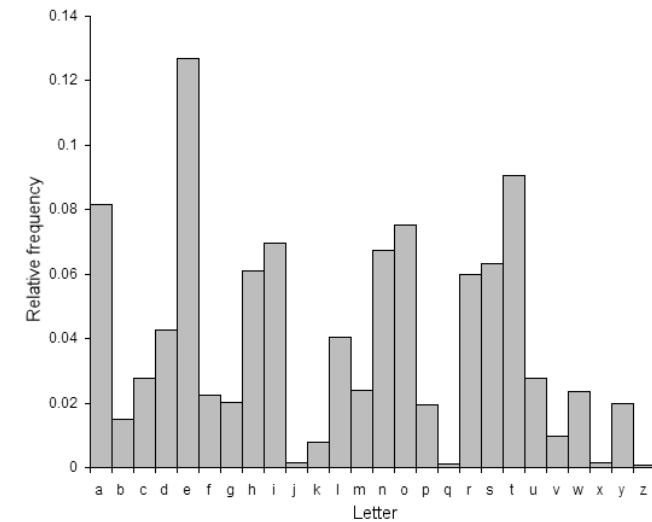


**statistical attacks**

$$\Phi_{\text{vigenere}}: \mathbb{Z}_{26} \rightarrow \mathbb{Z}_{26}$$

$$x \mapsto (x_i + k_{i \bmod t}) \bmod 26$$

message:            **A**TT**A**CK**A**TD**A**WN  
key:                **k**ey**k**ey**k**ey**k**ey  
ciphertext:        **K**XR**K**G**I**K**X**B**K**AL



# Vigenere cipher

$$k = (k_0, \dots, k_{t-1})$$

$$H_{vigenere} \simeq 4.7 \cdot t$$



**statistical attacks**



**key reuse, simple keys**

$$\varphi_{vigenere}: \mathbb{Z}_{26} \rightarrow \mathbb{Z}_{26}$$

$$x \mapsto (x_i + k_{i \bmod t}) \bmod 26$$

message:           ATTACKATDAWN  
key:                keykeykeykey  
ciphertext:        KXRKGIKXBKAL

COMPLETE VICTORY  
MANCHESTER BLUFF  
COME RETRIBUTION

# Vigenere cipher

$$k = (k_0, \dots, k_{t-1})$$

$$H_{vigenere} \approx 4.7 \cdot t$$



**statistical attacks**



**key reuse, simple keys**



**complex**

message:       ATTACKATDAWN  
key:            keykeykeykey  
ciphertext :   KXRKGIKXBKAL

A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z
B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A
C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B
D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C
E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D
F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E
G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F
H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G
I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H
J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I
K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J
L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K
M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L
N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M
O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N
P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O
Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P
R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	S
S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R
T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S
U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T
V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U
W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V
X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W
Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X
Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y

# Vernam cipher

$$k = (k_0, \dots, k_{n-1})$$

$$H_{\text{vernarn}} \simeq 4.7 \cdot n$$

$$\varphi_{\text{vernarn}}: \mathbb{Z}_{26} \rightarrow \mathbb{Z}_{26}$$

$$x_i \mapsto (x_i + k_{i \bmod n}) \bmod 26$$

message:	ATTACKATDAWN
key:	aezklwgrmali
ciphertext:	AXSKNGGKPOHV

# Vernam cipher

$$k = (k_0, \dots, k_{n-1})$$

$$H_{\text{vernarn}} \simeq 4.7 \cdot n$$

$$\varphi_{\text{vernarn}}: \mathbb{Z}_{26} \rightarrow \mathbb{Z}_{26}$$

$$x_i \mapsto (x_i + k_i \bmod n) \bmod 26$$

message:	ATTACKATDAWN
key:	aezklwgrmali
ciphertext:	AXSKNGGKPOHV

If  $k$  is uniformly distributed and independent on  $m$  this cipher is perfect, i.e.

$$P[M = m | C = c] = P[M = m]$$

# Vernam cipher



message:       ATTACKATDAWN  
key:            aezklwgrmali  
ciphertext :   KXRKGIKXBYAL

A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z
B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A
C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B
D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C
E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D
F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E
G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F
H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G
I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H
J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I
<b>K</b>	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J
L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K
M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L
N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M
O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N
P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O
Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P
R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	S
S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R
T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S
U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T
V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U
W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V
X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W
Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X
Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y







# Nomenclators



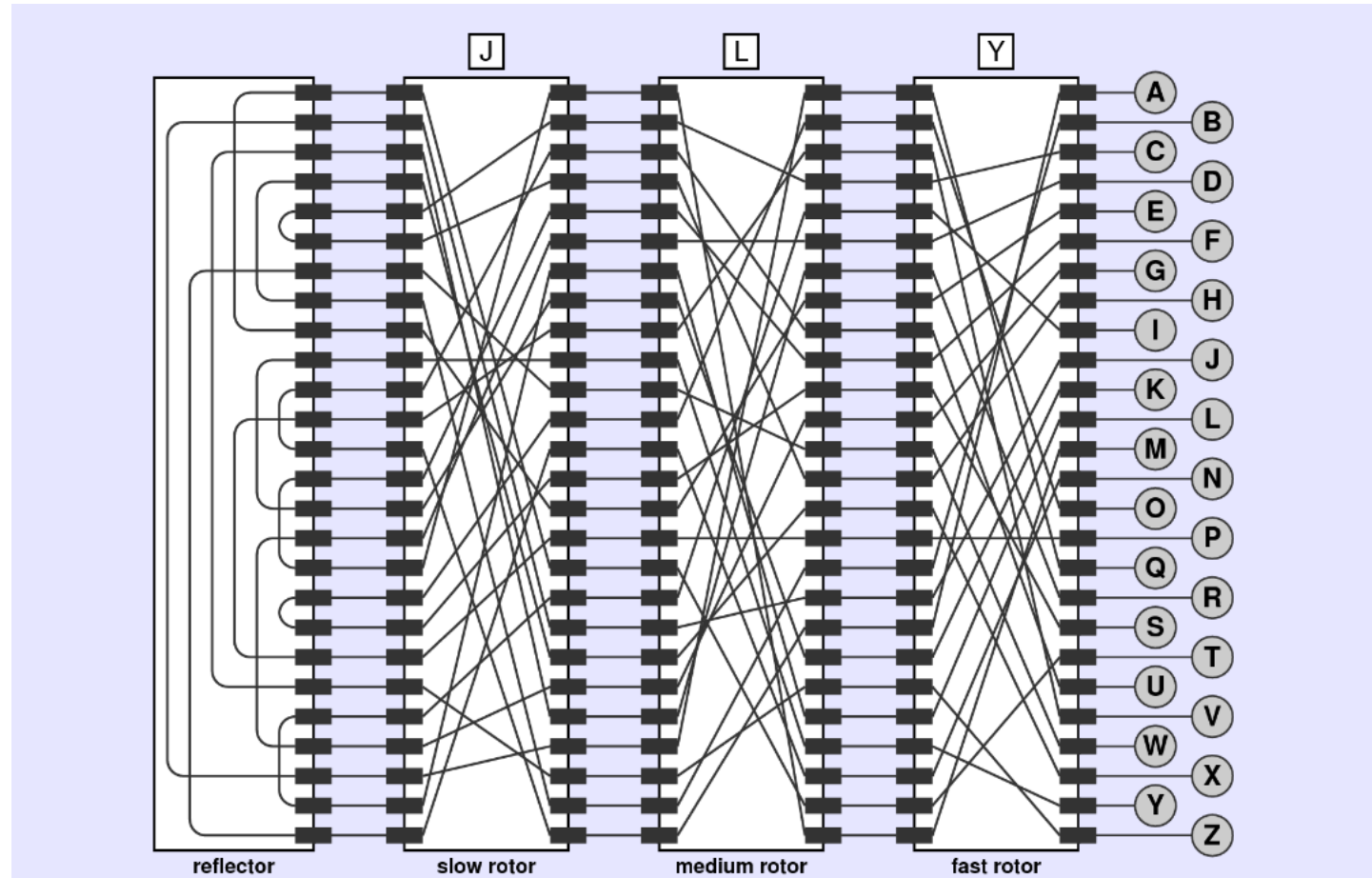
# Take aways

- Bad cryptography causes headaches/beheadings
- Short keys break
- Humans hate complex ciphers

**1.** Against entropy...

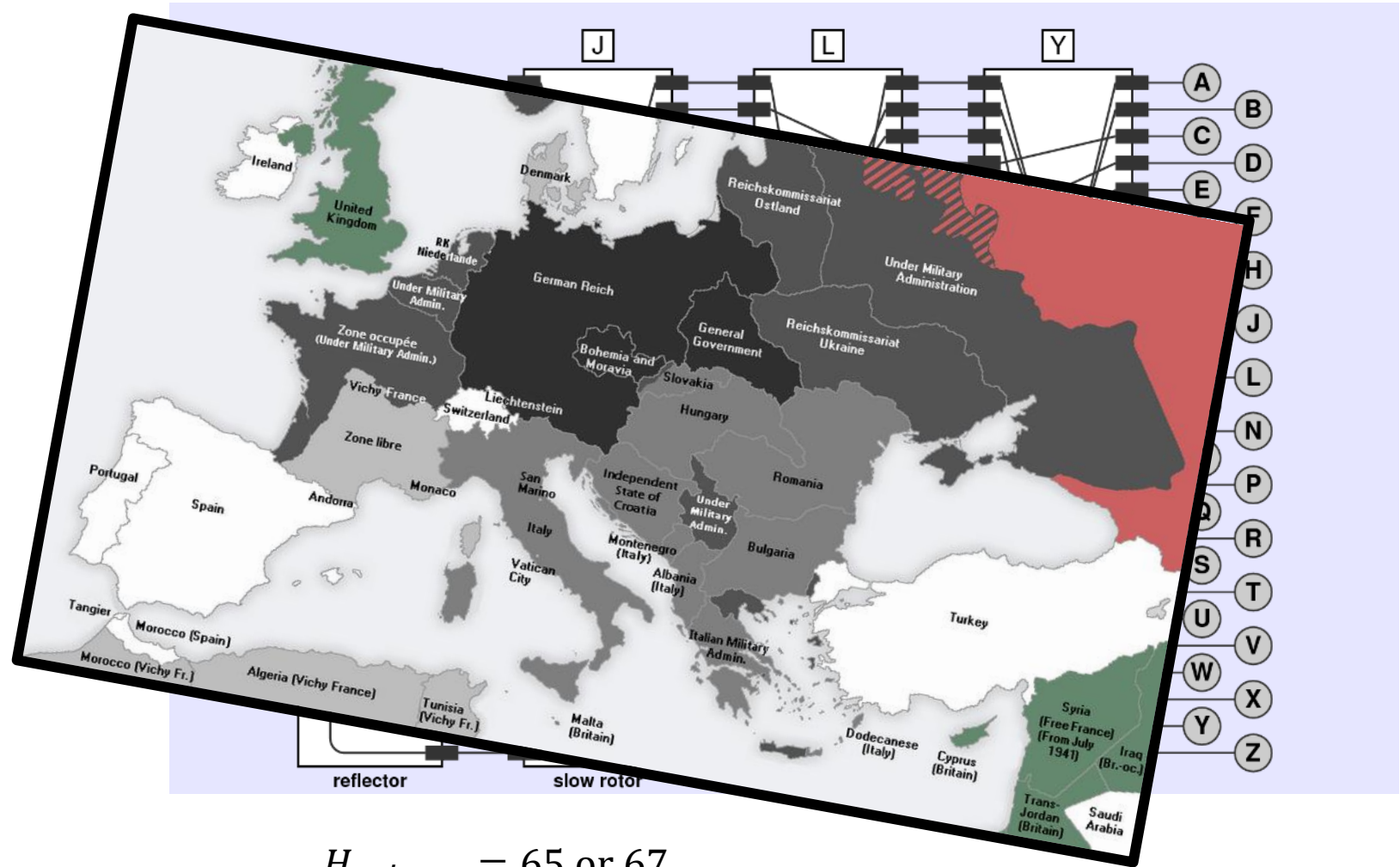
**2.** Machines themselves...

# Enigma cipher machine



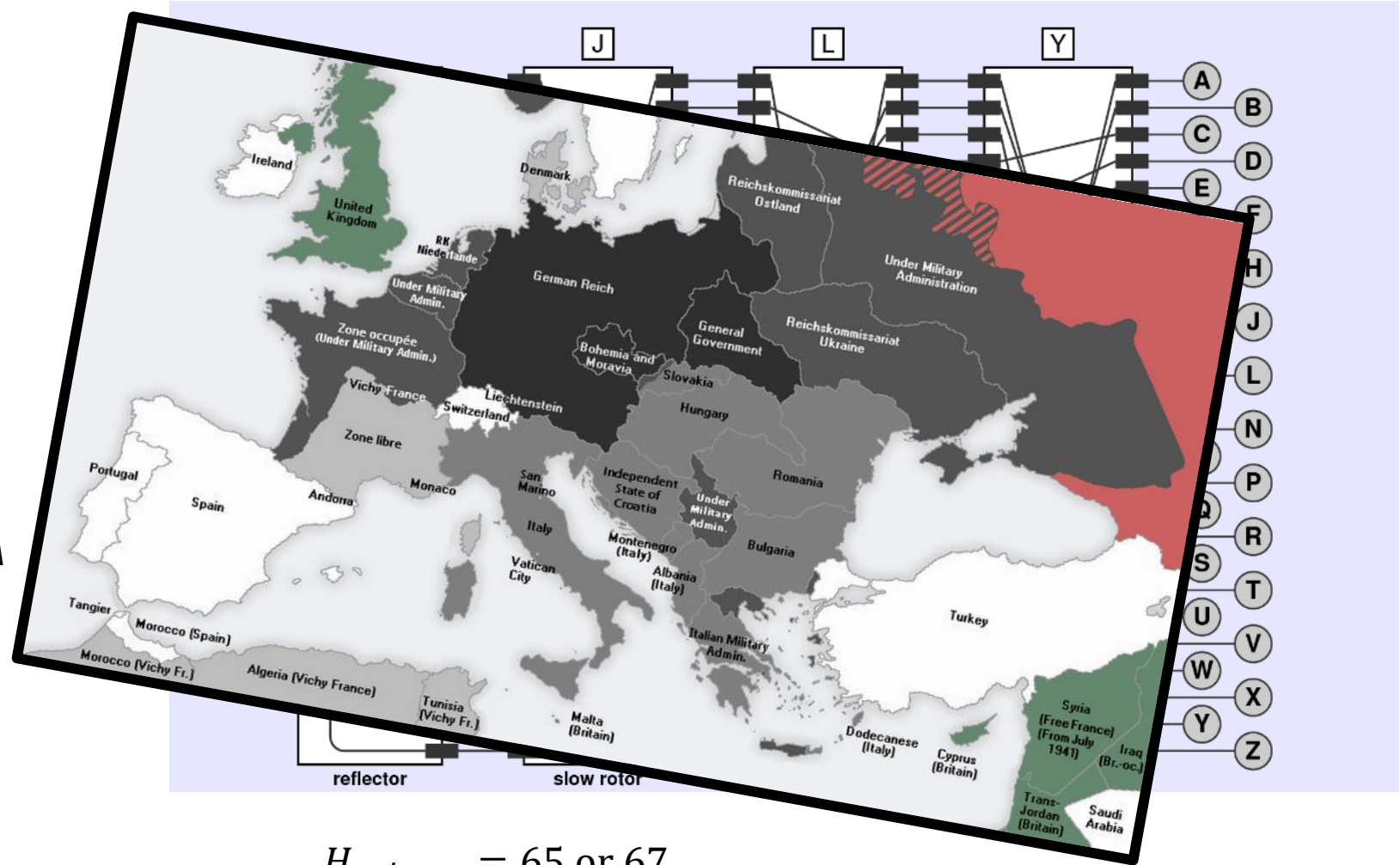
$$H_{\text{enigma}} = 65 \text{ or } 67$$

# Enigma cipher machine



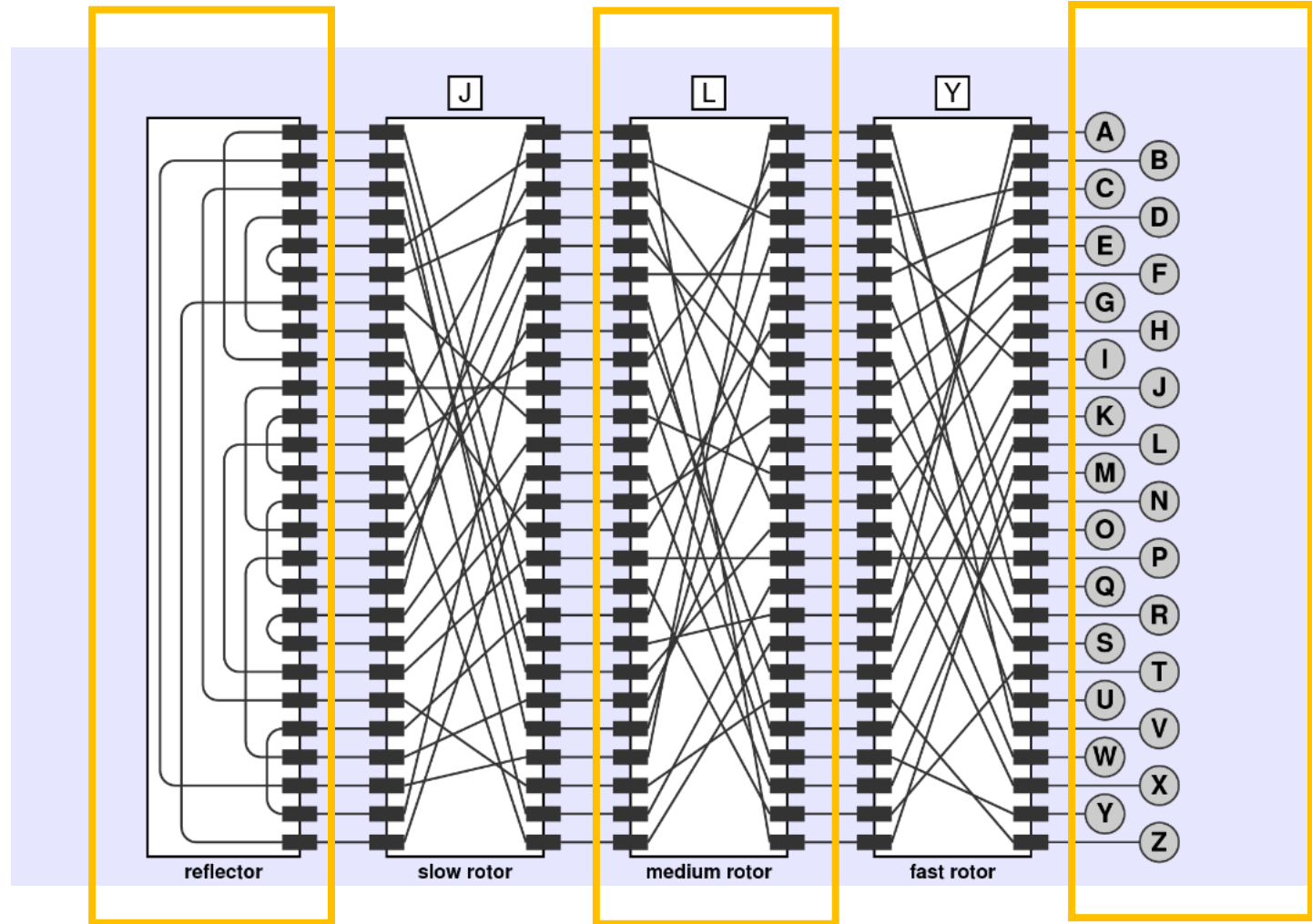
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# Enigma cipher machine



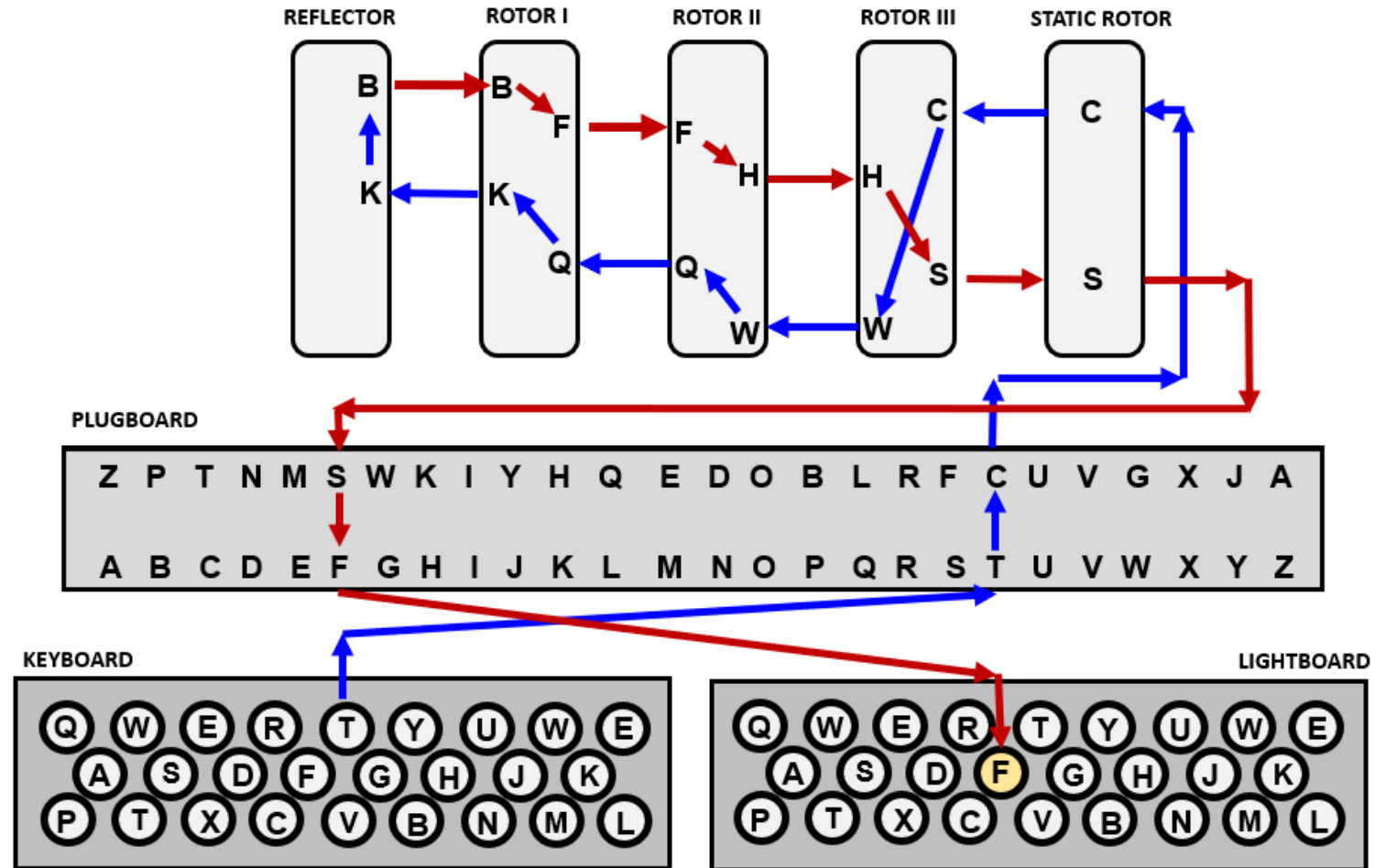
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# Enigma cipher machine



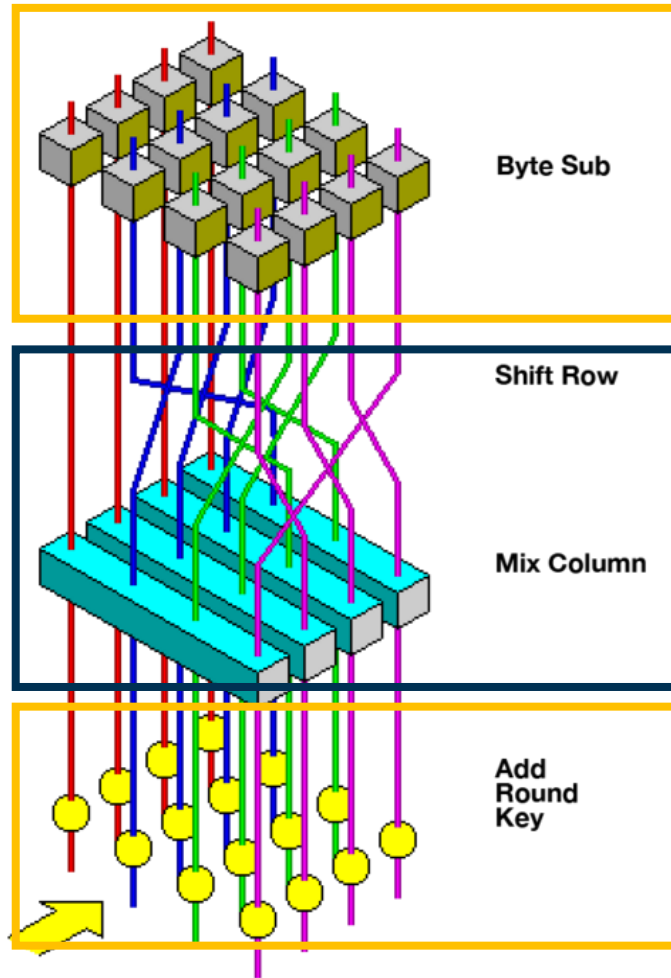
$$H_{\text{enigma}} = 65 \text{ or } 67$$

# Enigma cipher machine



$$H_{enigma} = 65 \text{ or } 67$$

# AES cipher



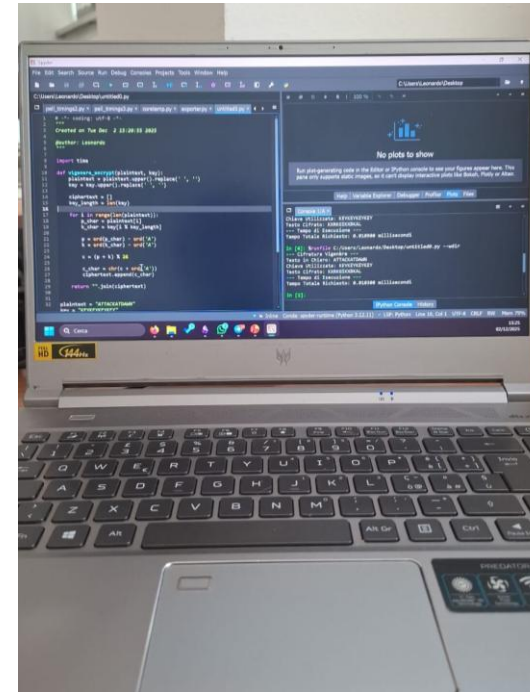
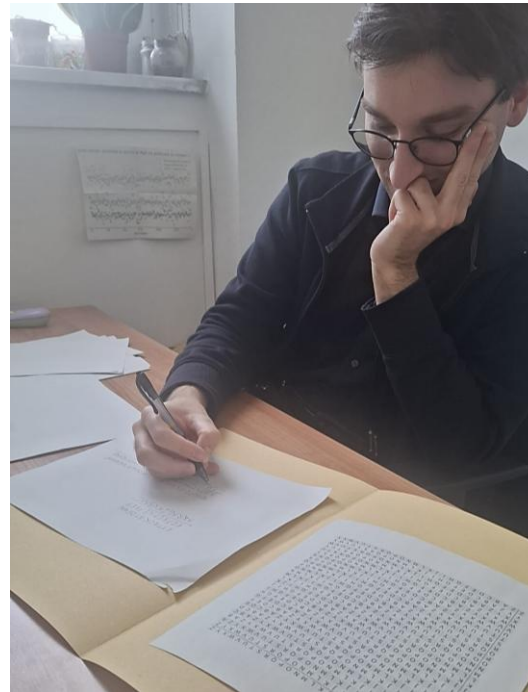
× 10 (or 12 or 14)

**Confusion:** each output character should depend on the key in a complex way.

**Diffusion:** each input character should influence many characters of the output, its statistical structure is dissipated.

$$H_{AES} = 128 \text{ or } 192 \text{ or } 256$$

# Good news



*Vigenere*  
*AES-128*

**me**

15 minutes

6 to 8 hours

**my laptop**

0.00001 seconds

0.00015 seconds

*\* benchmark on 128 characters*

# Bad news



# New fronteers

644

IEEE TRANSACTIONS ON INFORMATION THEORY, VOL. IT 22, NO. 6, NOVEMBER 1976

## New Directions in Cryptography

*Invited Paper*

WHITFIELD DIFFIE AND MARTIN E. HELLMAN, MEMBER, IEEE

*Abstract*—Two kinds of contemporary developments in cryptography are examined. Widening applications of teleprocessing have given rise to a need for new types of cryptographic systems, which minimize the need for secure key distribution channels and supply the equivalent of a written signature. This paper suggests ways to solve these currently open problems. It also discusses how the theories of communication and computation are beginning to provide the tools to solve cryptographic problems of long standing.

### I. INTRODUCTION

**W**E STAND TODAY on the brink of a revolution in

The best known cryptographic problem is that of privacy: preventing the unauthorized extraction of information from communications over an insecure channel. In order to use cryptography to insure privacy, however, it is currently necessary for the communicating parties to share a key which is known to no one else. This is done by sending the key in advance over some secure channel such as private courier or registered mail. A private conversation between two people with no prior acquaintance is a common occurrence in business, however, and it is unrealistic to expect initial business contacts to be postponed long enough for keys to be transmitted by some physical means. The cost and delay imposed by this key distribution

“We stand today on the brink of a revolution in cryptography. The development of cheap digital hardware has freed it from the design limitations of mechanical computing [...]. At the same time, theoretical developments [...] show promise of providing provably secure cryptosystems, changing this ancient art into a science.”

---

(W. Diffie, M. Hellman)

# New fronteers 1: paradigms

## PRIVATE KEY

# #

a very large  
secret prime  
number

a very large  
secret prime  
number



## PUBLIC KEY

#

the product of those  
two very large prime  
numbers used to make  
the private key, which  
is very, very hard to  
reverse back

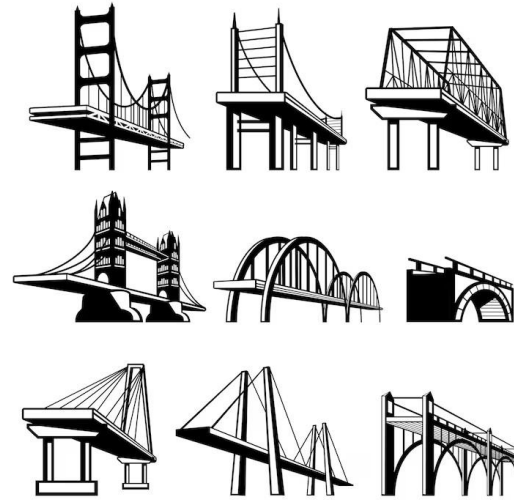


# New fronteers1: paradigms

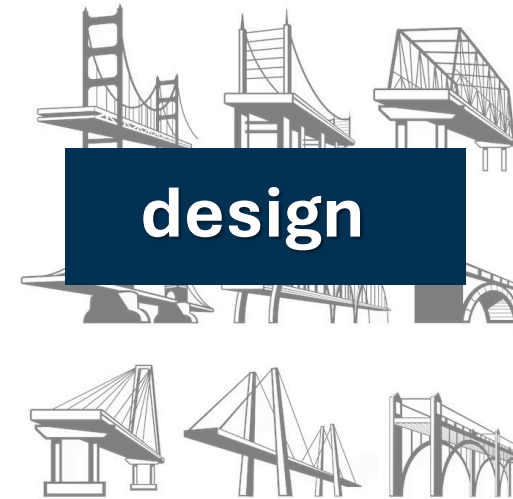
- Block ciphers
- Hash functions
- Stream ciphers
- Signature schemes
- Zero knowledge proofs
- Key encapsulation mechanisms
- ...



# New fronteers: 2, security



# New fronteers: 2, security



**INTERLUDE: What Cryptography is Not (p2)**



FANNY



ARTHUR



CORALINE



HIPPOLYTE



GWEN



AUGUSTIN



MAGGIE



ROGER



MARTIN



THEOPHILE



CAPUCINE



MAXIME



MELINA



ROBIN

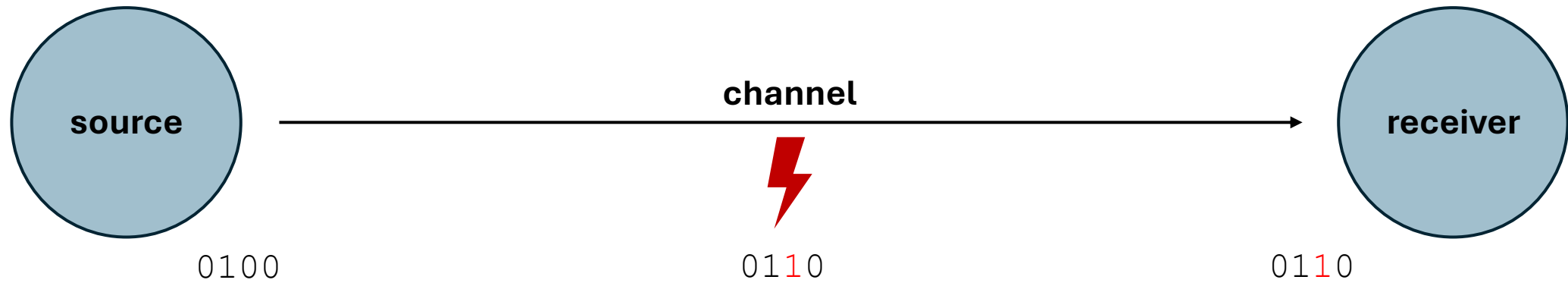


SANDRA

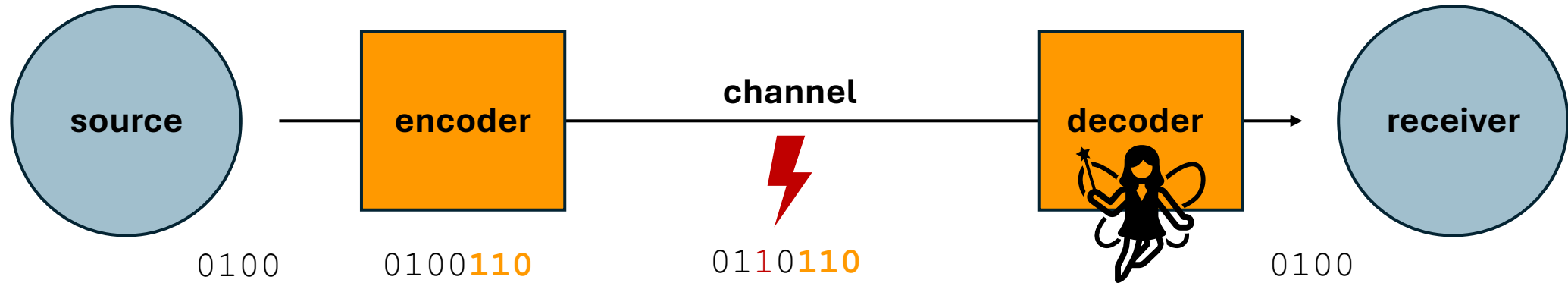


JUSTINE

# Interlude: coding theory



# Interlude: coding theory



$A$  alphabet  
 $k$  input length  
 $n$  code length

# Interlude: coding theory



Hackaday

## Cosmic Ray Flips Bit, Assists Mario 64 Speedrunner

The leading theory is that this bit flip was caused by a cosmic ray event, though the likelihood of such an event is exceedingly rare. It's...

17 feb 2021



BBC

## Bit flips: How cosmic rays grounded a fleet of aircraft

Radiation from space that led to more than 6000 Airbus aircraft needing emergency computer updates could become a growing problem.

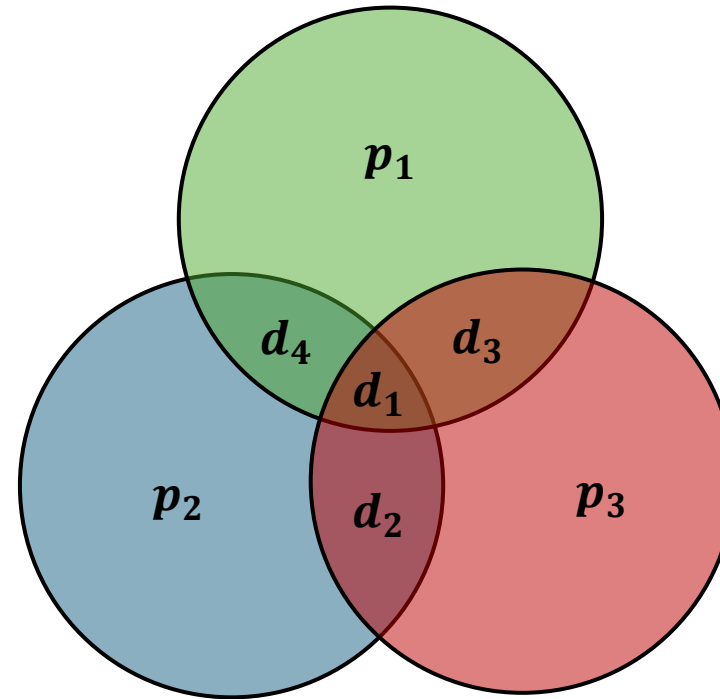
1 dic 2025



*“around one error per month per 256 MiB of RAM on a desktop computer”  
(IBM, 1996)*

# Interlude: coding theory

0	0000 <b>000</b>
1	0001 <b>011</b>
2	0010 <b>101</b>
3	0011 <b>110</b>
4	0100 <b>110</b>
5	0101 <b>101</b>
6	0110 <b>011</b>
7	0111 <b>000</b>
8	1000 <b>111</b>
9	1001 <b>100</b>
10	1010 <b>010</b>
11	1011 <b>001</b>
12	1100 <b>001</b>
13	1101 <b>010</b>
14	1110 <b>100</b>
15	1111 <b>111</b>



$$p_1 = d_1 + d_2 + d_4$$

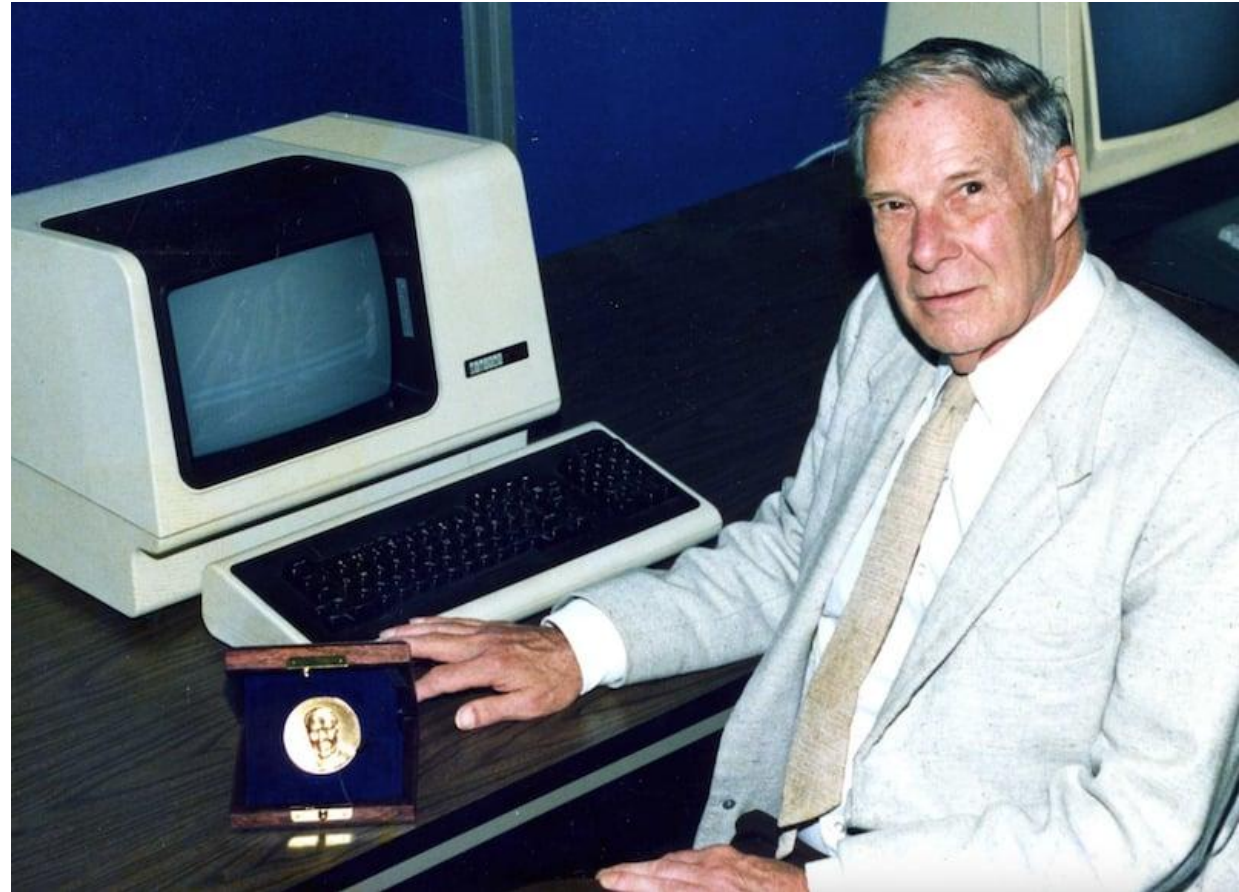
$$p_2 = d_1 + d_3 + d_4$$

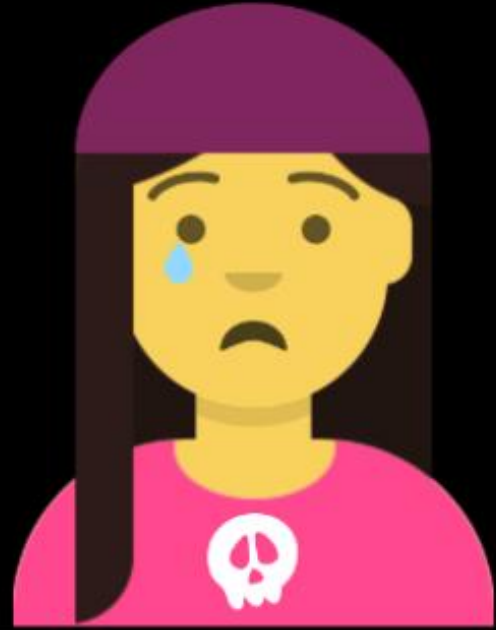
$$p_3 = d_2 + d_3 + d_4$$

# Interlude: coding theory

0	0000	000
1	0001	011
2	0010	101
3	0011	110
4	0100	110
5	0101	101
6	0110	011
7	0111	000
8	1000	111
9	1001	100
10	1010	010
11	1011	001
12	1100	001
13	1101	010
14	1110	100
15	1111	111

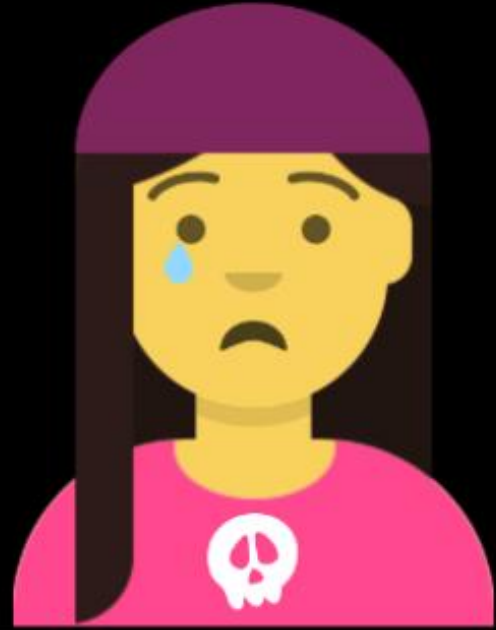
[7,4] Hamming code





**CAPUCINE**

**1010**

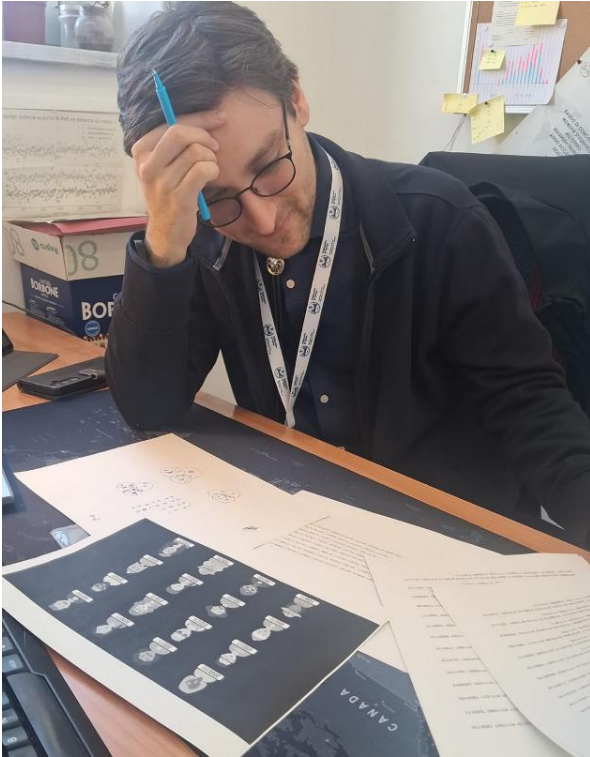


**CAPUCINE**

**1010101**



# Interlude: coding theory



*Me, decoding*



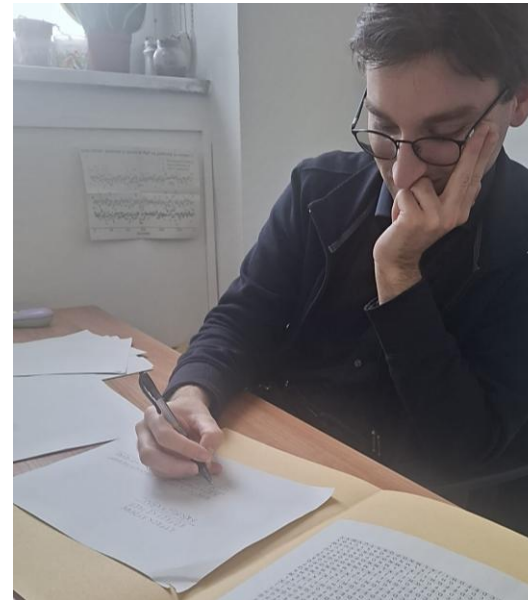
```
answer with 'y' for Yes or 'n' for No:  
1. Hat? 1  
2. Glasses? 0  
3. Is sad? 1  
4. Moustache? 0  
5. Drawing on shirt? 1  
6. Blue shirt/garment? 1  
7. Dark hair? 1  
Lie detected at question 6!  
Identified Person: CAPUCINE  
Do you want to play again? (y/n):
```

*My laptop, decoding*

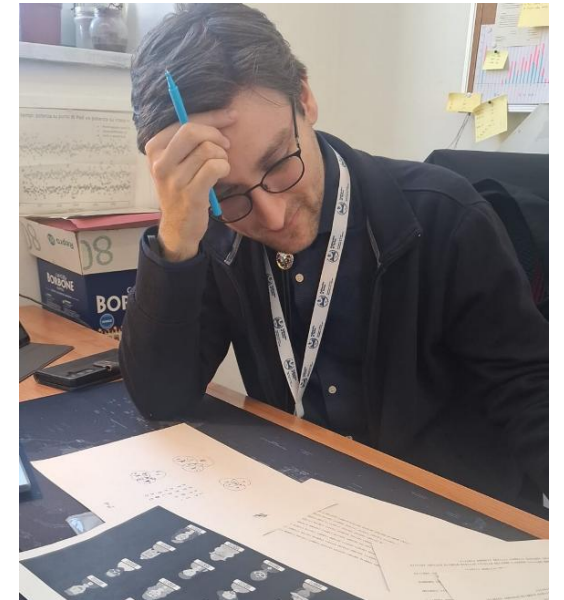
# Take aways

- Ciphers are now tailored for machines
- New directions: paradigms, security
- Cryptography  $\neq$  coding theory

- Mathematics is not stressful at all



*Encrypting (Dec. 2025)*



*Decoding (Feb. 2026)*

- 1. Against entropy...**
- 2. Machines themselves...**
- 3. Contend in vain?**

# Human error: Enigma



■ Simple keys

AAA, ABC, CIL

# Human error: Enigma



- Simple keys
- Repeated messages

«Wetter für die Nacht..»

# Human error: Enigma

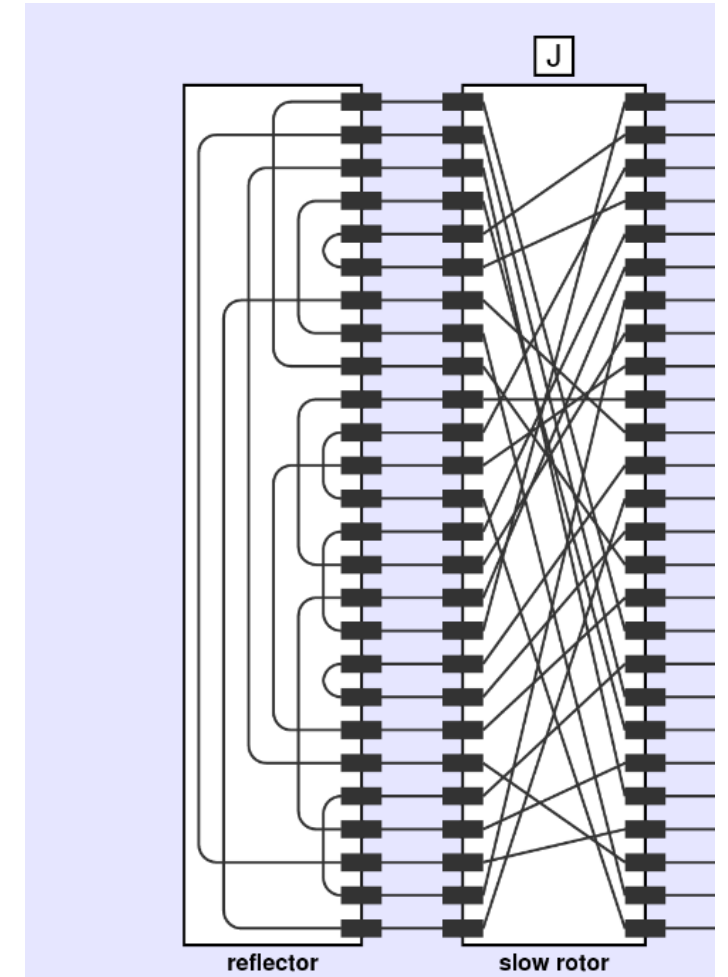


- Simple keys
- Repeated messages
- Short messages

# Human error: Enigma



- Simple keys
- Repeated messages
- Short messages
- Reflector (!)



# Human error: Sony PS3



secret key  $sk$



# Human error: Sony PS3



secret key *sk*



```
C5 B2 BF A1 A4 13 DD 16 F2 6D
31 C0 F2 ED 47 20 DC FB 06 70
```

# Government Backdoors 101

## **What?**

Any method to bypass security layers

# Government Backdoors 101

## What?

Any method to bypass security layers

## Why?

National security, law enforcement



*The four horsemen of infocalypse: drug-dealers, money-lauderers, terrorists, and pedophiles.*

# Government Backdoors 101

## What?

Any method to bypass security layers

## Why?

National security, law enforcement

## Are they ethical?

Who knows, probably not much



*The four horsemen of infocalypse: drug-dealers, money-lauderers, terrorists, and pedophiles.*

# Government Backdoors 101

## What?

Any method to bypass security layers

## Why?

National security, law enforcement

## Are they ethical?

Who knows, probably not much

## Are they safe?

No



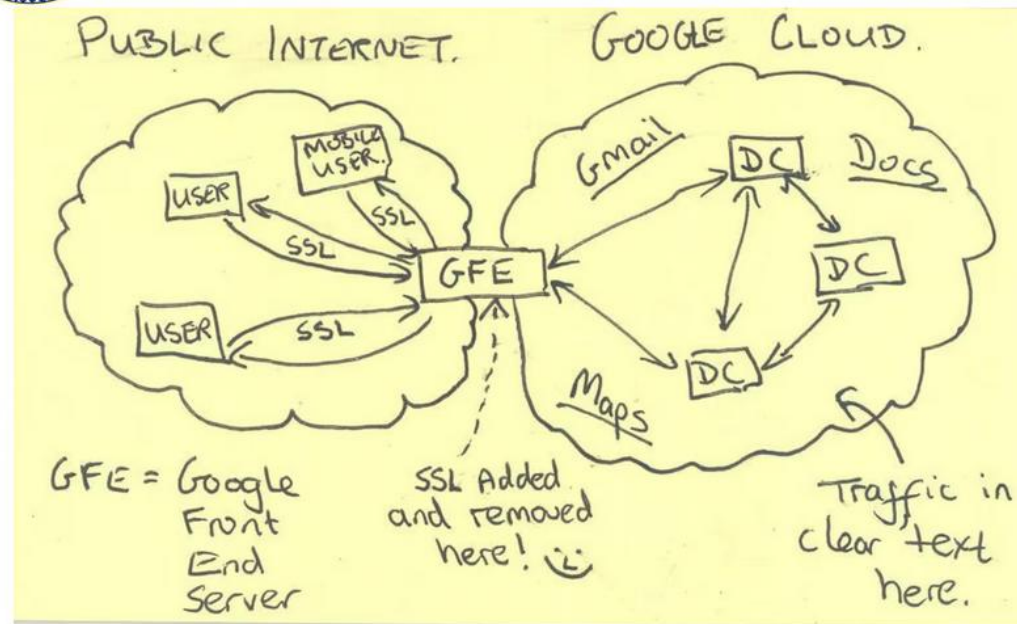
*The four horsemen of infocalypse: drug-dealers, money-launderers, terrorists, and pedophiles.*

# Government Backdoors 101

TOP SECRET//SI//NOFORN



## Current Efforts - Google



TOP SECRET//SI//NOFORN



**This will try to be as objective as possible.  
Still, it's hard to be objective when NSA slides look like this.**

# Backdoors: Operation Rubicon

1951: Boris Hagelin meets NSA



# Backdoors: Operation Rubicon

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1952: Boris Hagelin founds CryptoAG



# Backdoors: Operation Rubicon

1951: Boris Hagelin meets NSA

1952: Boris Hagelin founds CryptoAG

1970: CIA & BND buy CryptoAG via shell companies



=



+



*CIA + Bundesnachrichtendienst*

# Backdoors: Operation Rubicon

1951: Boris Hagelin meets NSA

1952: Boris Hagelin founds CryptoAG

1970: CIA & BND buy CryptoAG via shell companies



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+



# Backdoors: Operation Rubicon

1951: Boris Hagelin meets NSA

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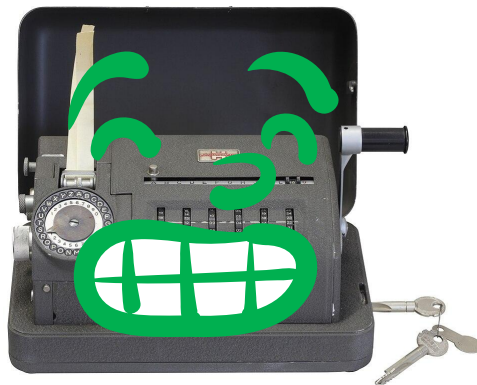
1970: CIA & BND buy CryptoAG via shell companies



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+



# Backdoors: Operation Rubicon

1951: Boris Hagelin meets NSA

1952: Boris Hagelin founds CryptoAG

1970: CIA & BND buy CryptoAG via shell companies

1991: IRAQ takes Hans Bühler hostage



=



+



# Backdoors: Operation Rubicon

1951: Boris Hagelin meets NSA

1952: Boris Hagelin founds CryptoAG

1970: CIA & BND buy CryptoAG via shell companies

1991: IRAQ takes Hans Bühler hostage

1994: Hans Bühler spills the beans



=



+



# Backdoors: Operation Rubicon

1973: US coup in Chile

1978: Egypt-Israel Camp David Accords

1982: Falklands war

1986: US attacks on Libia (Operation El Dorado Canyon)

1989: US invasion of Panama (Operation Just Cause)

+ *many, many domestic events*

# Backdoors: Operation Rubicon

## America

Argentina  
Brazil  
Chile  
Colombia  
Honduras  
Mexico  
Nicaragua  
Peru  
Uruguay  
Venezuela

## Europe

Austria  
Belgium  
Czechoslovakia  
Greece  
Hungary  
Ireland  
Italy  
Portugal  
Romania  
Spain  
Türkiye  
Vatican City  
Yugoslavia

## Africa

Algeria  
Angola  
Congo  
Egypt  
Gabon  
Ghana  
Guinea  
Ivory Coast  
Libya  
Mauritius  
Morocco  
Nigeria  
South Africa  
Sudan  
Tanzania  
Tunisia  
Zaire  
Zimbabwe

## Middle East

Iran  
Iraq  
Jordan  
Kuwait  
Lebanon  
Oman  
Qatar  
Saudi Arabia  
Syria  
U.A.E.

## Asia

Bangladesh  
Burma  
India  
Indonesia  
Japan  
Malaysia  
Pakistan  
Phillipines  
South Korea  
Thailand  
Vietnam

## Organisations

United Nations

*"The Intelligence Coup of the Century"*  
- CIA



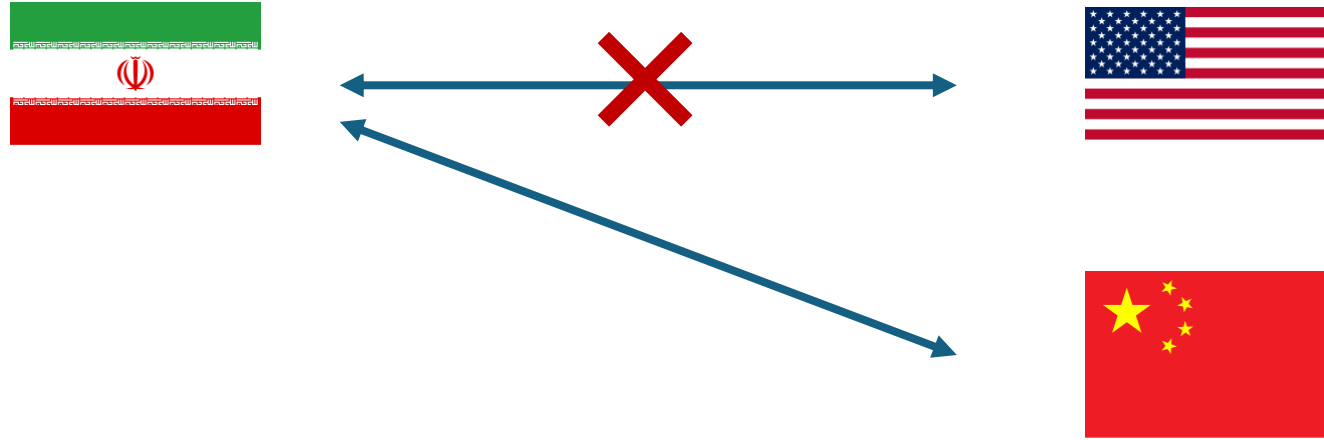
# Backdoors: Operation Rubicon



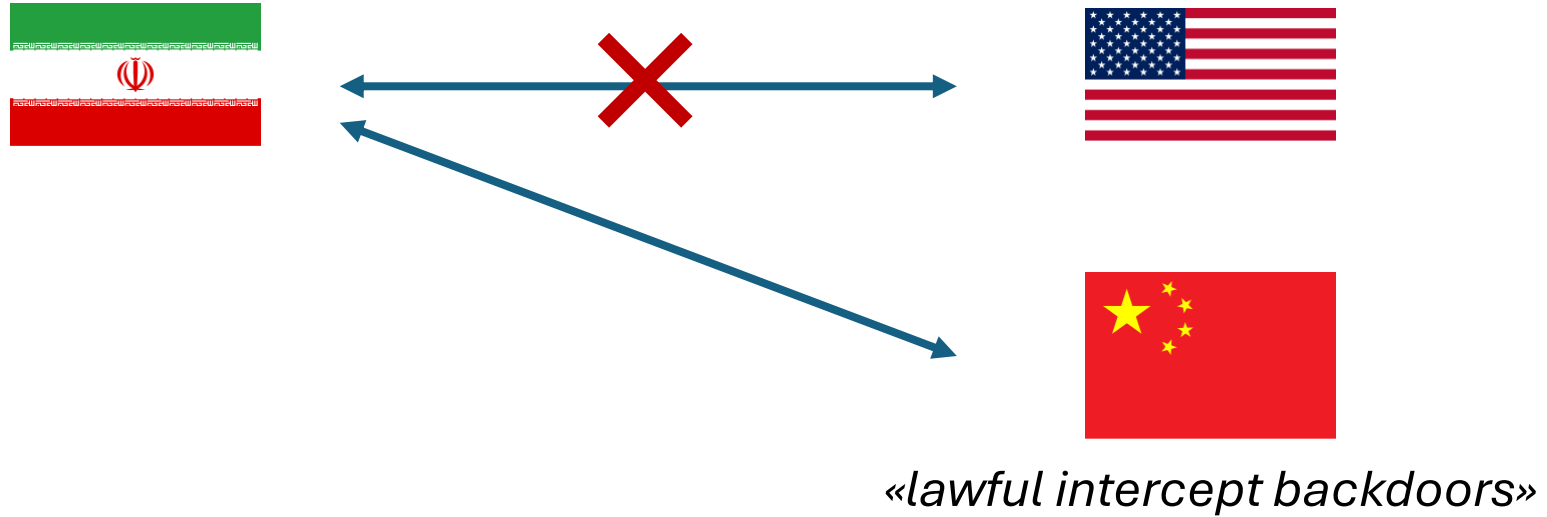
# Backdoors: Operation Rubicon

*Italian traffic was reportedly still being deciphered around 2001.*

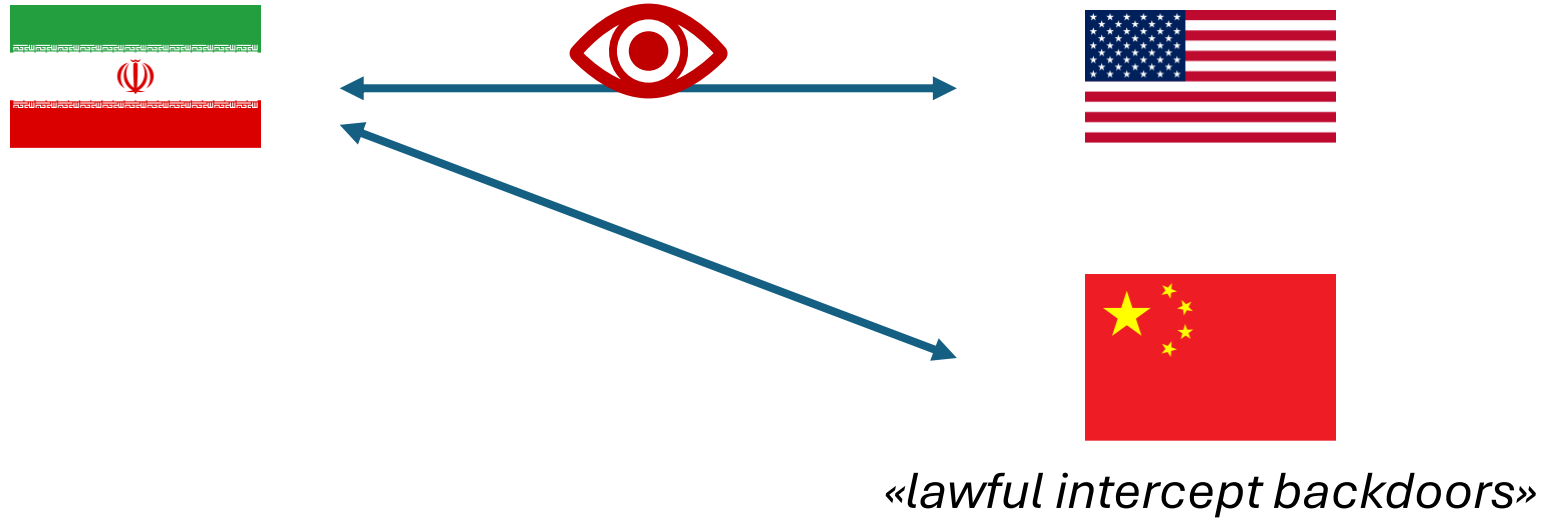
# Backdoors backfire: Operation Shotgiant



# Backdoors backfire: Operation Shotgiant



# Backdoors backfire: Operation Shotgiant



THN The Hacker News

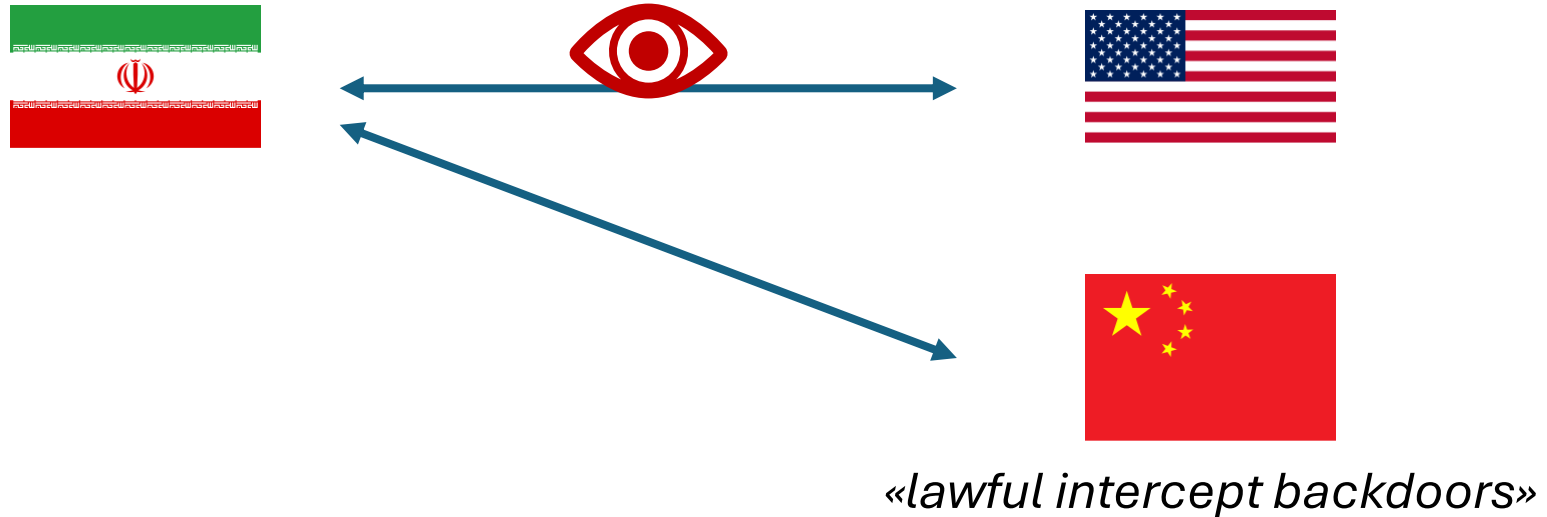
## NSA Hacked Servers of Chinese telecom Huawei, Stole Source Codes

Operation Shotgiant: NSA Hacked Chinese telecom Huawei and Stole Product Source Codes to write exploits.

23 mar 2014



# Backdoors backfire: Operation Shotgiant



*[...] this is both curious and puzzling. Have you ever seen a thief in the street who advertises [...] that he is a thief? Honestly speaking, I think what the U.S. has done here cannot be accepted as correct.*

- Hua Chunying, Foreign Ministry spokesperson

# CryptoWars™

## 1991: Phil Zimmerman & PGP

- Cryptography is a weapon, export prohibited
- Source code printed on books
- Case closed in 1996 thanks to public pressure

## 1999: Bernstein vs United States

- Publishing is protected by 1st emendament

## Conclusion: pyrrhic victory

- Cryptography can be exported, but...
- RSA export limited to weak keys



**US export regulations**

**Phil Zimmerman**

# Backdoors: NSA & DualEC

2006: NSA proposes a PRNG algorithm for adoption

2007: Shumow & Ferguson find something that *looks like a backdoor*

*“we're not saying NSA did this, but...”*

## The Main Point

- If an attacker knows  $d$  such that  $d*P = Q$  then they can easily compute  $e$  such that  $e*Q = P$  (invert mod group order)
- If an attacker knows  $e$  then they can determine a small number of possibilities for the internal state of the Dual Ec PRNG and predict future outputs.
- We do not know how the point  $Q$  was chosen, so we don't know if the algorithm designer knows  $d$  or  $e$ .

## Conclusion

- **WHAT WE ARE NOT SAYING:**  
NIST intentionally put a back door in this PRNG
- **WHAT WE ARE SAYING:**  
The prediction resistance of this PRNG (as presented in NIST SP800-90) is dependent on solving one instance of the elliptic curve discrete log problem.  
(And we do not know if the algorithm designer knew this before hand.)

# Backdoors: NSA & DualEC

2006: NSA proposes a PRNG algorithm for adoption

2007: Shumow & Ferguson find something that *looks like a backdoor*

*“we're not saying NSA did this, but...”*

2013: Post-Snowden, RSA Security admits getting USD 10Mln from NSA to adopt it

# Backdoors: NSA & DualEC

*“If you can’t beat them, trick them!”*  
*- NSA, probably*



## Backdoors: NSA & DualEC

## Backdoors: NSA & DES

*“If you can’t beat them, trick them!”*  
*- NSA, probably*



**Backdoors: NSA & DualEC**

**Backdoors: NSA & DES**

**Backdoors: NSA & AES**

*“If you can’t beat them, trick them!”  
- NSA, probably*



## Backdoors: NSA & post-quantum? (probably not)

CRYPTOGRAPHY

# **A very unscientific guide to the security of various PQC algorithms**



By Sophie Schmieg



December 13, 2025



No Comments

# Backdoors: Chat Control



*“Proposal for a REGULATION OF THE EUROPEAN PARLIAMENT AND OF THE COUNCIL laying down rules to prevent and combat child sexual abuse”*


Either

- Break/weaken E2E ← not a great idea
- Client-side scanning ← technically a spyware

# Backdoors: Chat Control




*“Proposal for a REGULATION OF THE EUROPEAN PARLIAMENT AND OF THE COUNCIL laying down rules to prevent and combat child sexual abuse”*

 Patrick Breyer

## End of “Chat Control”: EU Parliament Stops Mass Surveillance in Voting Thriller – Paving the Way for Genuine Child Protection!

The controversial mass surveillance of private messages in Europe is coming to an end. After the European Parliament had already rejected...

2 settimane fa

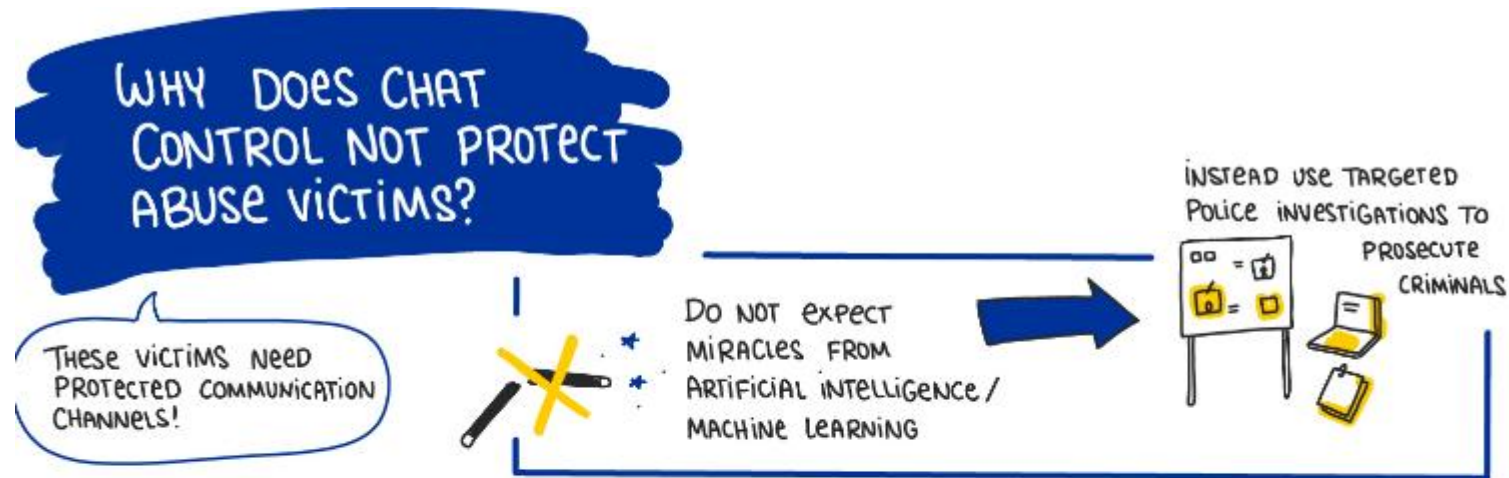


For: 307. Against: 306. Abstentions: 24. In total, 637 MEPs voted. 81 MEPs didn't vote.

# Backdoors: Chat Control



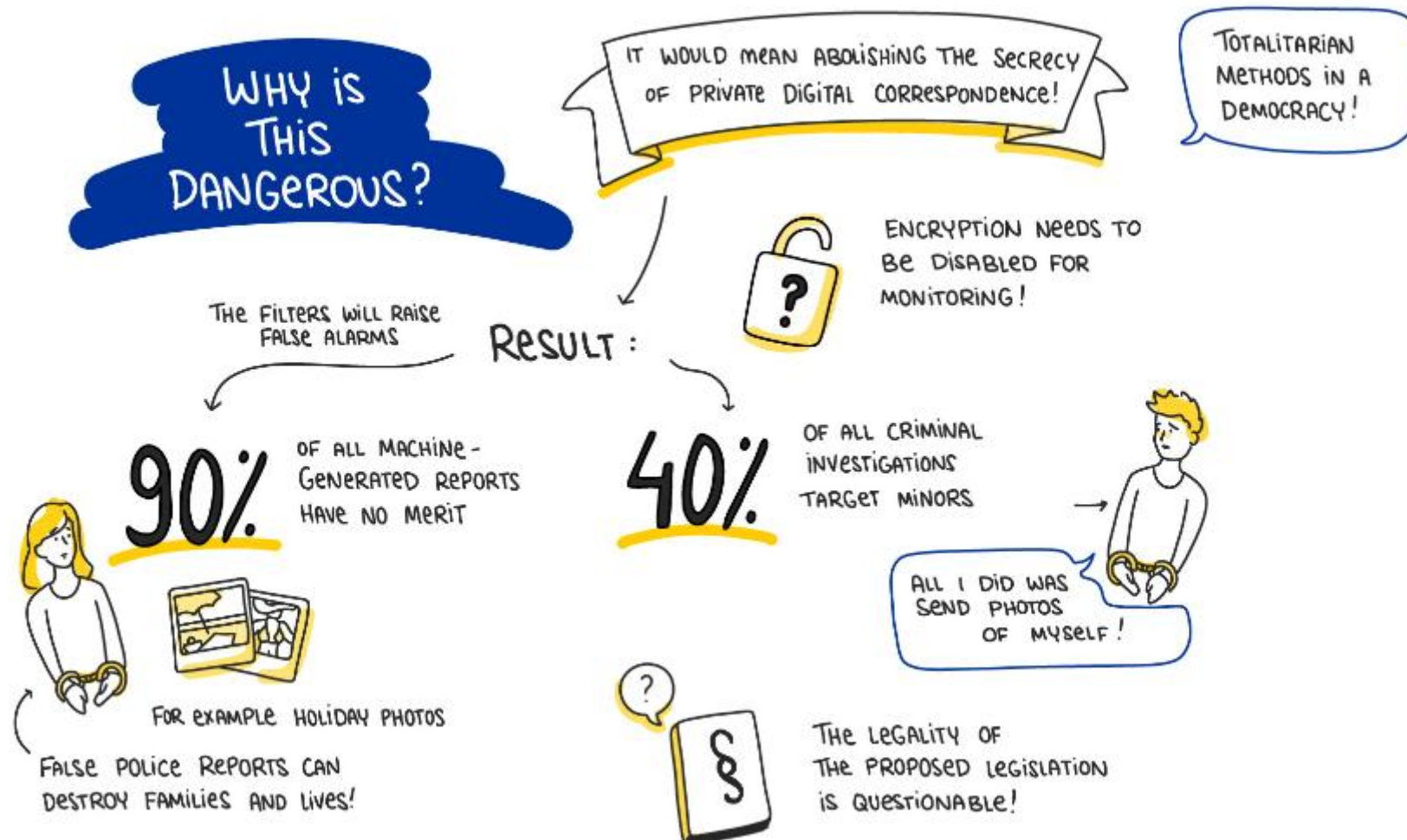
“Proposal for a REGULATION OF THE EUROPEAN PARLIAMENT AND OF THE COUNCIL laying down rules to prevent and combat child sexual abuse”



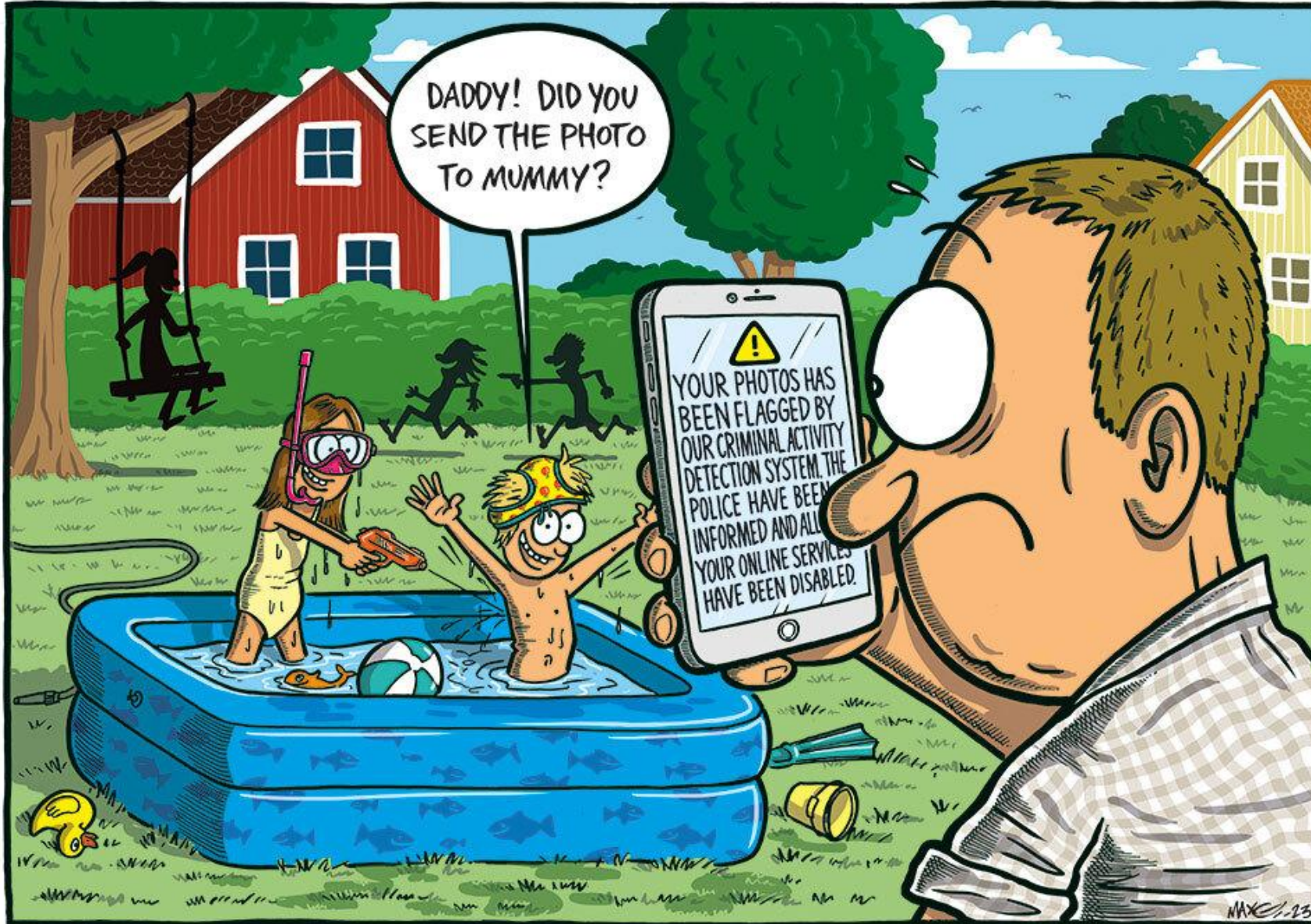
# Backdoors: Chat Control



“Proposal for a REGULATION OF THE EUROPEAN PARLIAMENT AND OF THE COUNCIL laying down rules to prevent and combat child sexual abuse”



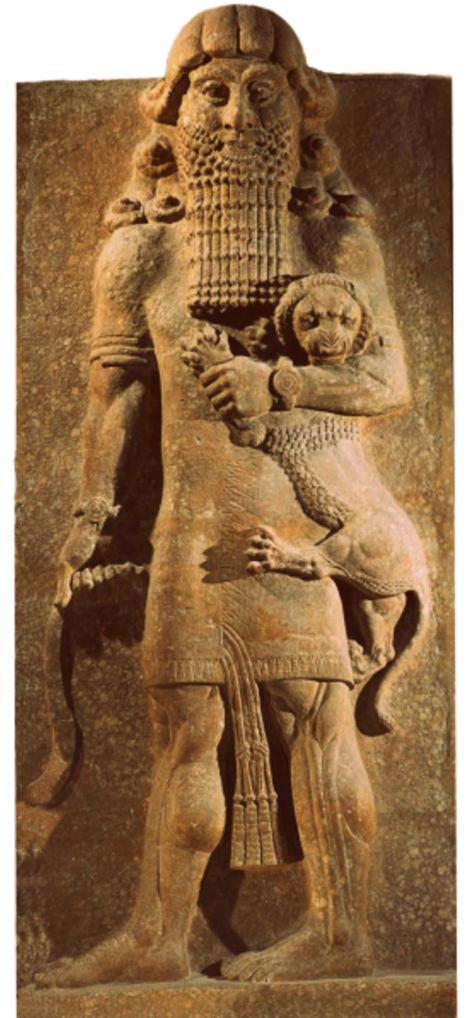
# Backdoors: Chat Control



# Take aways

- Humans are bad at following protocols. Really bad.
- Tension between surveillance and privacy
- How secure can cryptography be?

*“Get in, fair reader, into the depths of Cryptography.  
Inspect the theory, view its foundations.  
Is not the very core made of fine mathematics?”*  
- **Gilgamesh**



# HisTo26 HISTORY OF CRYPTOGRAPHY IN TORINO

TORINO, ITALY  
APRIL 17TH 2026

A seminar on the history of cryptography, in the historical city of Turin.

**HisTo26** explores the evolution of secret writing, from Classical and Renaissance ciphers to World War II cryptanalysis and the cold war. Topics include: evolution of cryptography, historical ciphers, analysis of documents and cases of study.

**Organisers:** Paolo Bonavoglia (DeCifris), Leonardo Errati (PoliTo).

## Invited speakers:



**Dr. Dermot Turing**  
Visiting Fellow  
Kellow College, Oxford



**Prof. Joachim Rosenthal**  
Professor of Mathematics  
University of Zurich



**Flavio Atzeni**  
Chief Warrant Officer  
Italian CC (retired)



**Dr. Francesco Cosimato**  
Brigadier General  
Italian Army (retired)

With the contribution of BitPolito and Disma.  
<http://sites.google.com/view/histo26/home>



Politecnico  
di Torino



BitPolito



Politecnico  
di Torino

Department  
of Mathematical Sciences  
"G. L. Lagrange"

